

CS 4100 Homework 03

Due Sunday 2/19 at midnight (1 minute after 11:59 pm) in Gradescope (with a grace period of 6 hours)

You may submit the homework up to 24 hours late (with the same grace period) with NO penalty (just this once :-)

You must submit the homework in Gradescope as a zip file containing **two files**:

- The `.ipynb` file (be sure to Kernel \rightarrow Restart and Run All before you submit); and
- A `.pdf` file of the notebook.

For best results obtaining a clean PDF file on the Mac, select File \rightarrow Print Review from the Jupyter window, then choose File \rightarrow Print in your browser and then Save as PDF. Something similar should be possible on a Windows machine.

PLEASE collapse the Utility Code for Problems 1 - 3 before creating the PDF so that we don't have to search for your solutions to the problems.

All homeworks will be scored with a maximum of 100 points; if point values are not given for individual problems, then all problems will be counted equally.

Utility Code for Problems 1 - 3

The cells below, until the heading **Problem One**, implements a prover in first-order logic; don't change anything in these cells without discussing with Prof Snyder first.

If you have enabled code folding (an excellent idea), you can collapse this utility code so that it does not display, but is still runnable.

In [1]:

```
# Pretty Printing

def literal2String(L):
    if(type(L)==int):
        return ('-' if (L<0) else '') + chr(ord('A')+abs(L)-1)
    elif(L[0]!='not'):
        return 'not('+term2String(L[1])+')'
    else:
        return term2String(L)

def clause2String(C):
    if(len(C)==0):
        return '{} '
    return '{ '+(','+'.join([literal2String(L) for L in C]))+ '}'

def clause2Sequent(C):
    if(len(C)==0):
        return '<='
    RHS = [ A[1] for A in C if A[0]!='not' ]
    LHS = [ A for A in C if A[0]!='not' ]
    if(len(RHS)==0 and len(LHS)==1):
        return term2String(LHS[0])
    return ('+'.join([literal2String(L) for L in LHS])) + '<=' + ('+'.join([literal2String(L) for L

def clauseSet2String(A):
    return '{ '+(','+'.join([clause2String(C) for C in A]))+ '}'

def clauseList2String(A):
    return '[' '+(','+'.join([clause2String(C) for C in A]))+ ']'

def term2String(t):
    if(len(t) == 1):
        return t[0]
    else:
        return t[0] + '(' + ('+'.join([term2String(s) for s in t[1:]])+')'
```

```

def substitution2String(s):
    return '{ '+(',' .join([x + ' : ' + term2String(t) for (x,t) in s.items()])+ ' }'

def pprint(X1,X2=None):
    pprintAux(X1)
    if(X2!=None):
        print('\t',end='')
        pprintAux(X2)
    print()

def pprintAux(X):
    #print('pprintAux:',X)
    if(type(X)==int or type(X)==tuple):           # X is literal
        print(literal2String(X),end='')
    elif(type(X)==set):                           # X is a clause
        print(clause2String(X),end='')
    elif(type(X)==list):                          # X is a clause list
        print(clauseList2String(X),end='')
    elif(type(X)==dict):                          # X is a substitution
        D = { k:term2String(X[k]) for k in X.keys() }
        print(D,end='')
    else:
        print('Error in pprintAux!')

#test

clauseSet2String([ [1], [-1,2], [-2,-1,3] ])

s = ('f',('x',), ('g',('a',)))
t = ('f',('h',('z',)),('z',))
#clauseSet2String([ [('P',s)], [('not',('Q',s,t))] ])

A = ('P',('f',('x',)))
NA = ('not',('P',('f',('x',))))
#literal2String(A)
clause2String( { ('P',('f',('x',))),('not',('Q',('h',('x',)))),('not',('P',('f',('x',)))),'Q',('g',('x',)) }
clauseList2String( [{ A,NA },{ A,NA }])
clause2Sequent( { ('P',('f',('x',))),('not',('Q',('h',('x',)))),('not',('P',('f',('x',)))),'Q',('g',('x',)) }

# pprint(-2)
# pprint({3,-2})
# pprint([ {1},{-1,2}, {-2-1,3} ])
# pprint([ {1}, {}, {-2} ])
# pprint(('P',('f',('x',))))
# pprint({'not',('Q',('g',('x',),('a',)))),'P',('f',('x',))})
# pprint([ {'P',('g',('f',('x',)),('f',('b',)))}, ('not',('P',('f',('x',)))) ] )
# pprint( { 'x': ('a',), 'y': ('f',('g',('x',))) } )

```

Out[1]: 'P(f(x)), Q(g(x)) <= Q(h(x)), P(f(x))'

In [2]:

```

# parser

def getTokens(s):
    lo = hi = 0
    tokens = []
    s = s.replace(' ', '') # remove blanks
    while(hi < len(s)):
        c = s[hi]
        if(c.isalnum()):
            hi += 1
        elif(lo==hi): # found ( ,
            tokens += [ s[lo:(hi+1)] ]
            lo = hi = hi+1
        else: # found string
            tokens += [ s[lo:hi] ]
            lo = hi
    return tokens

def getTerm(t):
    return getTermAux(getTokens(t+'$'),0)[0]

def getTermAux(ts,k): # get next term starting at index k, return encoding and next index after
    # must start with string
    if(ts[k][0].isalnum()):

```

```

    if(ts[k+1] == '('):
        (tlt,next) = getTermList(ts,k+2)
        return (tuple([ts[k]] + tlt),next)
    else:
        #elif(ts[k+1] == ')'):
        return ((ts[k],),k+1)
else:
    return "Error2"

def getTermList(ts,k): # get next term list starting at index k, return encoding and next index after
    TL = []
    while(True):
        (t,k) = getTermAux(ts,k)
        TL.append(t)
        if(ts[k]==' '):
            return (TL,k+1)
        elif(ts[k]==','):
            k += 1

# parseClause({'not(Grandmother(karen,f(x,a)))', 'P(x)', 'not(Q(a,y))'})

```

In [3]:

```

# SLR parser

def reportError(s,i, follow):
    print('Parsing error: expecting\n')
    print('\t',end='')
    for c in follow:
        print(c+' ',end='')
    print('\n\nhere:\n')
    print('\t'+s[:-1]+' \n\t'+(' '*i)+'^')

def parse(s,trace=False):
    s += '$'
    stack = [1]
    i=0
    while(True):
        if trace:
            print(s[i],stack)
        state = stack[-1] # top state on stack
        if(state == 0):
            if(s[i] == '$'):
                return stack[1]
            else:
                reportError(s,i,['<end of expression>'])
                return
        elif(state in [1,3,5]):
            if(s[i].isalnum()):
                stack += [s[i],6]
                i += 1
            else:
                reportError(s,i,['<letter or number>'])
                return
        elif(state == 2):
            if(s[i] == '('):
                stack += [s[i],3]
                i += 1
            elif(s[i] in ')','$'): # reduce by 1: E -> S
                exp = stack[-2]
                stack.pop()
                stack.pop()
                if(stack[-1]==1):
                    stack += [(exp,),0]
                elif(stack[-1] in [3,5]):
                    stack += [(exp,),4]
            else:
                reportError(s,i,['(',')'],'<letter or number>', '<end of expression>'])
                return
        elif(state == 4):
            if(s[i] == ','):
                stack += [s[i],5]
                i += 1
            elif(s[i] == ')'): # reduce by 5: EL -> E
                exp = stack[-2]

```

```

        stack.pop()
        stack.pop()
        if(stack[-1]==3):
            stack += [[exp],9]
        elif(stack[-1]==5):
            stack += [[exp],8]
    else:
        reportError(s,i,['(',')'])
        return
elif(state == 6):
    if(s[i].isalnum()):
        stack += [s[i],6]
        i += 1
    elif(s[i] in '(),$'): # reduce by 3: S -> let
        let = stack[-2]
        stack.pop()
        stack.pop()
        if(stack[-1] in [1,3,5]):
            stack += [let,2]
        elif(stack[-1] == 6):
            stack += [let,7]
    else:
        reportError(s,i,['(',')','<letter or number>','<end of expression>'])
        return
elif(state == 7):
    if(s[i] in '(),$'): # reduce by 4: S -> let S
        exp = stack[-4]+stack[-2]
        stack.pop()
        stack.pop()
        stack.pop()
        stack.pop()
        if(stack[-1] in [1,3,5]):
            stack += [exp,2]
        elif(stack[-1] == 6):
            stack += [exp,7]
    else:
        reportError(s,i,['(',')','(',')','<end of expression>'])
        return
elif(state == 8):
    if(s[i] in '$'): # reduce by 6: EL -> E,EL
        exp = [stack[-6]] + stack[-2]
        stack.pop()
        stack.pop()
        stack.pop()
        stack.pop()
        stack.pop()
        stack.pop()
        if(stack[-1]==3):
            stack += [exp,9]
        elif(stack[-1]==5):
            stack += [exp,8]
    else:
        reportError(s,i,['()'])
        return
elif(state == 9):
    if(s[i]==')'):
        stack += [s[i],10]
        i += 1
    else:
        reportError(s,i,['()'])
        return
elif(state == 10):
    if(s[i] in '$'): # reduce by 2: E -> S(EL)
        exp = tuple([stack[-8]] + stack[-4])
        stack.pop(); stack.pop()
        stack.pop(); stack.pop()
        stack.pop(); stack.pop()
        stack.pop(); stack.pop()
        if(stack[-1]==1):
            stack += [exp,0]
        elif(stack[-1] in [3,5]):
            stack += [exp,4]
    else:
        reportError(s,i,['(',')','(',')','<end of expression>'])
        return

```

```

def parseClause(S):      # S is a list of strings representing atomic formulae
    return { parse(t) for t in S }

def parseClauseList(L):
    return [ parseClause(C) for C in L ]

s = 'not(P(father(x234,a),y))'
# s = 'father(x)'
# s = 'x4'
# print(s)
# parse(s)

# parseClause({'not(Grandmother(karen,f(x,a)))', 'P(x)', 'not(Q(a,y))'})
# a = parseClauseList({'not(Grandmother(karen,f(x,a)))', 'P(x)', 'not(Q(a,y))'}, {'not(Brother(karen,f(x
# pprint(a)

```

In [4]:

```

# Unification and substitutions

def isVar(t):
    return len(t)==1 and t[0][0] in ['u','v','w','x','y','z'] # any string starting with u,v,w,x, y,

def isConst(t):
    return len(t)==1 and not (t[0][0] in ['u','v','w','x','y','z'])

def getVars(t,vs=set()):
    if type(t)==set:
        for k in t:
            vs = vs.union(getVars(k,vs))
        return vs
    elif isVar(t):
        return {t[0]}
    elif isConst(t):
        return set()
    else:
        for k in range(1,len(t)):
            vs = vs.union(getVars(t[k],vs))
        return vs

def occursIn(v,t):
    if isVar(t):
        return ((v,) == t)
    elif isConst(t):
        return False
    else:
        return any( [ occursIn(v,s) for s in t[1:] ] )

def applySubst(t,subst):
    if type(t)==set:
        return { applySubst(c,subst) for c in t }
    elif isVar(t) and t[0] in subst.keys():
        return subst[t[0]]
    elif isConst(t):
        return t
    else:
        return tuple([t[0]] + [ applySubst(t[k],subst) for k in range(1,len(t)) ] )

def composeSubst(sub1,sub2):
    s = { v : applySubst(t,sub2) for (v,t) in sub1.items() }
    s.update(sub2)
    return s

def unify(s,t,subst={}):
    s = applySubst(s,subst)
    t = applySubst(t,subst)
    if s==t:
        return subst
    if isVar(t):
        (s,t) = (t,s)
    if isVar(s):
        if occursIn(s[0],t):
            return None
        else:
            return composeSubst(subst,{s[0]:t})
    if s[0] != t[0] or len(s)!=len(t):

```

```

    return None
for k in range(1,len(s)):
    subst = unify(s[k],t[k],subst)
    if(subst==None):
        return None
return subst

# s = ('x',)
# t = ('a',)
# s = ('f',('g',('x',)),)
# t = ('f',('x',))
# s = ('f',('x',), ('g',('a',)))
# t = ('f',('h',('z',)),('y',))

# print(term2String(s))
# print(term2String(t))
# sub = unify(s,t)
# pprint(sub)
# print(term2String(applySubst(s,sub)))
# print(term2String(applySubst(t,sub)))

```

In [5]:

```

# renaming of variables

seed = 0

def reseed():
    global seed
    seed = 0

def getNewVariable():
    global seed
    seed += 1
    return ('x' + str(seed-1),)

def rename(t):
    rs = { v : getNewVariable() for v in getVars(t)}
    return applySubst(t,rs)

# C = { ('P',('x',),('y',)), ('P',('y',('f',('x',)))), ('not',('Q',('x',),('z',))) }
# pprint(C)
# pprint(rename(C))

```

In [6]:

```

# Resolution rule

def resolveFOL(C1,C2):
    C1 = rename(C1)
    C2 = rename(C2)
    R = []
    for L1 in C1:
        for L2 in C2:
            if(L1[0]=='not' and L2!='not'): # complementary literals
                usub = unify(L1[1],L2)
                if(usub != None):
                    T1 = set(C1)
                    T1.remove(L1)
                    T2 = set(C2)
                    T2.remove(L2)
                    R.append( applySubst(T1.union(T2),usub) )
            elif(L2[0]=='not' and L1!='not'): # complementary literals
                usub = unify(L2[1],L1)
                if(usub != None):
                    T2 = set(C2)
                    T2.remove(L2)
                    T1 = set(C1)
                    T1.remove(L1)
                    R.append( applySubst(T1.union(T2),usub) )
    return R

def resolveAllFOL(A,C):
    R = [ resolveFOL(C1,C) for C1 in A ]
    return [C for CL in R for C in CL]

```

In [7]:

```

# prover for FOL
# due to difficulty with renaming, tautology and duplicate checking is not done
# only refinement is to order the queue by size

def proveFOL(KB,SOS,limit=30,trace=False):
    Queue = SOS
    count = 0 # count the number of pops off the queue
    while(len(Queue) > 0):
        if(trace):
            print('Queue:', clauseList2String(Queue))
            # form all resolvents with front of queue,
            # check if empty clause is generated, else add to end of queue
            C = Queue.pop()
            count += 1
            for C1 in resolveAllFOL(KB + Queue,C):
                if(len(C1) == 0):
                    print("\nUnsatisfiable! (" ,count,"step(s) executed ")
                    return
                else:
                    Queue = [C1] + Queue # BFS
                    Queue.sort(reverse=True, key=(lambda x: len(x)))
            if(count >= limit):
                print("\nEmpty clause not found after", limit,"steps!")
                return
        print("\nSatisfiable! (" ,count,"step(s) executed ")

# test
# KB = [ { ('not',('P',('f',('x',))))}, ('Q',('x',)) },{ ('not',('Q',('x',))) } ]
# NQ = [ {('P',('x',)) } ]

# print('KB: ',end='')
# pprint(KB)
# print('NQ: ',end='')
# pprint(NQ)
# print()
# proveFOL(KB,NQ)

```

Problem One (30 pts)

Part A

This part is just to make you familiar with the workflow for the FOL prover. It includes a parser for input, which will make writing formulae easier, since you can specify atomic formula as strings, which are then converted into the internal representation (as suggested in the textbook on p.) used by the parser.

You will not need to write any of the prover, and all of the clauses you should have to write will be Horn clauses (although the prover will work with non-Horn clauses).

The code for printing out the KB and SOS is written in "Prolog" style for ease of reading.

Please view the "walk-through" video for more explanation of how to enter clauses and run the prover.

For the following code, do the following:

- Replace `john` with `ketaki` in the knowledge base; and
- Replace the SOS with one which answers the following question: "Does Ketaki like peanuts?"

In [8]:

```

KB1a= [{ 'Food(apple)' },
        { 'Food(vegetables)' },
        { 'Eats(anil,peanuts)' },
        { 'Alive(anil)' },
        { 'not(Eats(y,z))', 'Food(z)', 'Killed(y)' },
        { 'not(Food(x))', 'Likes(john,x)' },
        { 'not(Eats(anil,w))', 'Eats(harry,w)' }
      ]

# parse to convert strings to encoding used by the prover

```

```

KB1a = parseClauseList(KB1a)

SOS1a = parseClauseList([ { 'not(Food(x))' } ])

print('KB1a: [')
for C in KB1a:
    print('\t',end='')
    print('clause2Sequent(C)')
print(']\n')
print('SOS1a: ',end='')
pprint(SOS1a)
print()

proveFOL(KB1a,SOS1a,trace=True)

```

```

KB1a: [
  Food(apple)
  Food(vegetables)
  Eats(anil,peanuts)
  Alive(anil)
  Food(z), Killed(y) <= Eats(y,z)
  Likes(john,x) <= Food(x)
  Eats(harry,w) <= Eats(anil,w)
]

SOS1a: [ { not(Food(x)) } ]

Queue: [ { not(Food(x)) } ]

Unsatisfiable! ( 1 step(s) executed )

```

Solution Part A

In []:

Part B

In this part you will define a knowledge base KB1b, which will be used in the rest of the problem to test various formulae to see if they are consequences of the knowledge base.

Encode the following knowledge base as shown in the book, defining the following predicates directly from the illustration:

```

Female(_)
Male(_)
Child(?,?,_)

```

As in the textbook, be sure to distinguish the two Oscars!

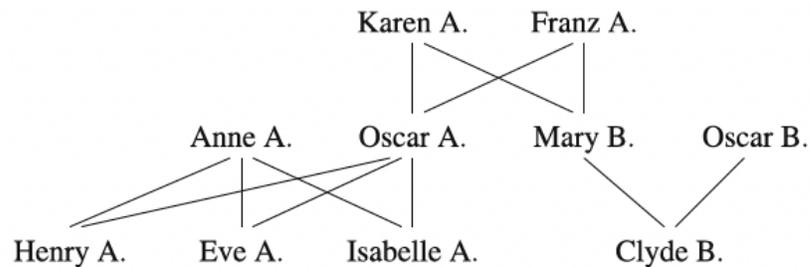


Fig. 3.1 A family tree. The edges going from Clyde B. upward to Mary B. and Oscar B. represent the element (Clyde B., Mary B., Oscar B.) as a child relationship

Next, define the following predicates:

```

Son(x,y)      -- x is the son of y
Daughter(x,y) -- x is the daughter of y
Father(x,y)   -- x is the father of y
Mother(x,y)   -- x is the mother of y
Parent(x,y)   -- x is a parent of y
Grandfather(x,y) -- x is the grandfather of y
Grandmother(x,y) -- x is the grandmother of y

```

Note the following carefully:

- You may NOT define any other predicates than these for the problems which follow.
- You may only use letters (upper or lower) and numbers in the symbols; variables start with u, v, w, x, y, or z.
- As proof of your solution for this part, simply parse and then print out the knowledge base as we did in Part A; but be sure to call your clause set `KB1b` !

In []:

Part C

Run the prover to answer the following question:

Is Karen a grandmother?

In []:

Part D

Run the prover to answer the following question:

Does Eve have a brother?

In []:

Part E

Run the prover to answer the following question:

Does Franz have both a grandson and a granddaughter?

(This may take a long time, so try setting the limit higher. My solution took 1175 steps.)

In []:

Problem Two (5 pts)

Suppose we want to write clauses for the predicates `Brother` and `Sister` as shown in the code cell below.

In [14]:

```

BS = [    {'Brother(x,y)', 'not(Child(x,z,w))', 'not(Child(y,z,w))', 'not(Male(x))' },
          {'Sister(x,y)', 'not(Child(x,z,w))', 'not(Child(y,z,w))', 'not(Female(x))' }
        ]

BS = parseClauseList(BS)

print('BS: [')
for C in BS:
    print('\t',end='')
    print clause2Sequent(C)
print('    ]\n')

```

```
BS: [
  Brother(x,y) <= Child(x,z,w), Male(x), Child(y,z,w)
  Sister(x,y) <= Child(x,z,w), Female(x), Child(y,z,w)
]
```

Part A

Unfortunately, this does not work as expected. For this problem, add these clauses to KB1b and try (each separately) to prove:

- Eve has a brother
- Henry has a brother
- Clyde has a brother

In []:

Part B

Explain why you think this is happening (you can't see the bindings for the variables, but you could look at the trace, or substitute names in place of variables to see if you can understand the issue.

Do you see an easy solution to this problem?

Answer Part B:

Problem Three (5 pts)

Part A

Run the prover from Problem Two to try to prove the following two assertions, one of which, in real life, would have to be true:

Clyde has a sister.

Clyde does not have a sister.

In []:

Part B

Describe what you see as the problem here, and suggest possible solutions (there may not be any really satisfactory ones, but think about it).

(Logically, this is a complex issue; you might want to Google "Closed-world assumption" and "Negation as failure" for background.)

Problem Four: The Eight-Puzzle (60 pts)

The Eight-Puzzle is described in the textbook and in lecture (and more thoroughly in the walk-through video for this homework). It is an interesting example of a non-trivial problem whose search space is large (it has $9! = 362,880$ distinct states) but not so large that you can't explore it completely in a reasonable amount of time.

In this problem we will develop a solver for this puzzle, and explore various search strategies for solving it.

Don't skip reading and thoroughly understanding the background material here before launching into Part A of the problem.

Utility code examples for the 8-Puzzle

In order to make the solver as efficient as possible we want to use two data structures which are more efficient than the naive approach of just storing everything in a list:

- A heap is a data structure for priority queues which has an $O(\log N)$ cost for adding and removing elements;
- A dictionary (which uses a heap) is an excellent Python data structure for storing paths between the states (boards) for recovering the exact structure of the shortest path.

We now demonstrate how these can be used for this problem.

In [20]:

```
# example of min heap for priority queue
# heapq works on a list in place

import heapq

# initializing list to be used as priority queue
PQ = []

print("The initialized heap is :",PQ)

# using heappush() to push elements into heap; the array PQ is changed in place

# if adding a tuple, it uses first element as the key
heapq.heappush(PQ, (4, 'A'))
heapq.heappush(PQ, (5, 'B'))
heapq.heappush(PQ, (2, 'C'))
heapq.heappush(PQ, (7, 'D'))

print("The heap after insertions is :",PQ,"\n")

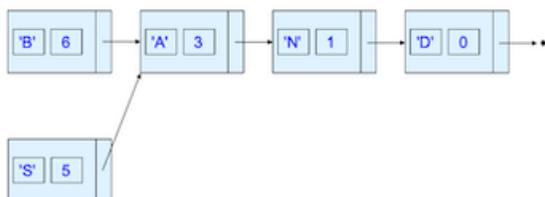
# pop the elements off in ascending order
# heappop(...) changes the list PQ in place and returns a minimal element
while(PQ):
    print("Smallest element:",heapq.heappop(PQ))
    print("The heap is now:",PQ)
```

```
The initialized heap is : []
The heap after insertions is : [(2, 'C'), (5, 'B'), (4, 'A'), (7, 'D')]
```

```
Smallest element: (2, 'C')
The heap is now: [(4, 'A'), (5, 'B'), (7, 'D')]
Smallest element: (4, 'A')
The heap is now: [(5, 'B'), (7, 'D')]
Smallest element: (5, 'B')
The heap is now: [(7, 'D')]
Smallest element: (7, 'D')
The heap is now: []
```

For storing predecessor lists for paths through the search space, what we essentially need is linked lists which store data (in our case, a nested tuple encoding the board), the heuristic function $f(x)$, and a link to the predecessor (or None, for the start state).

For simplicity, assume that we want to store paths through the capital letters 'A' .. 'Z' (with no duplicate letters), for example:



We can think of the letters as "states" through which we are searching, and we would like to implement a Closed list of states we have already processed, with the path back to the start state 'D'.

To implement this with a dictionary, we use the state as a key and a tuple to store the heuristic value and the previous state; you can quickly find any state and its value, change the value of a state, and trace back through the path, as shown in the next cell.

WARNING: You will need to use tuples (explained below) instead of letters for the states; you may not use lists, as they can not be hashed.

In [21]:

```

from collections import defaultdict

Closed = defaultdict((lambda x: None))

Closed['D'] = (0, None)           # f('D') = 0, and None represents Null
Closed['N'] = (1, 'D')           # Note carefully that 1 is f('N'), not f('D')
Closed['A'] = (3, 'N')
Closed['B'] = (6, 'A')
Closed['S'] = (5, 'A')

# test

print("f('B') =", Closed['B'][0])
print("pred('B') =", Closed['B'][1])
print()

# modify value of 'B'

new_value = 7
Closed['B'] = (new_value, Closed['B'][1])
print("f('B') =", Closed['B'][0])
print("pred('B') =", Closed['B'][1])
print()

# print out the path in order, using a loop, just as an example

# NOTE: in the problem you will need to print the path in reverse,
# which can be done by a recursive function

state = 'S'
while state:           # while not equal to None
    print(state)
    state = Closed[state][1]

```

```

f('B') = 6
pred('B') = A

```

```

f('B') = 7
pred('B') = A

```

```

S
A
N
D

```

Part A (25 pts)

The next cell contains some useful utility functions and the standard goal state/board from the textbook. The cell after contains stubs for functions you should write to implement the function `getChildren` which calculates the possible moves from a given board.

Please check with Prof Snyder if you wish to do something different.

In [22]:

```

from collections import defaultdict

Closed = defaultdict((lambda x: None))

# print out a simple representation of the state (board) with a blank line after

def printBoard(b):
    for r in range(3):
        for c in range(3):
            print(str(b[r][c]) + " ", end='')
        print()
    print()

```

```

# useful utility functions -- only work for 3x3 boards!

# convert nested tuple representation to nested lists
def tuples2Lists(b):
    return [list(rw) for rw in b]

# convert nested list representation to nested tuples
def lists2Tuples(b):
    return ((b[0][0],b[0][1],b[0][2]),
            (b[1][0],b[1][1],b[1][2]),
            (b[2][0],b[2][1],b[2][2]) )

# standard goal state

goalBoard = ((0, 1, 2), (3, 4, 5), (6, 7, 8))

printBoard(goalBoard)

b = tuples2Lists(goalBoard)
print(b)
gb = lists2Tuples(b)
print(gb)

```

```

0 1 2
3 4 5
6 7 8

```

```

[[0, 1, 2], [3, 4, 5], [6, 7, 8]]
((0, 1, 2), (3, 4, 5), (6, 7, 8))

```

In [23]:

```

# TODO:

# a "tile" is a number 0 - 8 on the board

# return the row in board b where tile t occurs
def row(t,b):
    pass # your code here

# return the column in board b where tile t occurs
def col(t,b):
    pass # your code here

# Heuristics from thetext book

# this is to do trials where no heuristic h is used
def noHeuristic(b1,b2):
    return 0

# count how many tiles are not in the correct position
def tilesOutOfPlace(b1,b2):
    pass # your code here

# sum the distances to move each tile to correct position if no other tiles existed
def manhattanDistance(b1,b2):
    pass # your code here

# return a list of boards that can be reached in one movement of tile 0
def getChildren(b):
    pass # your code here

# tests

btest1 = ((3,1,2),(4,5,0),(6,7,8))
btest2 = ((3,7,4),(8,0,2),(5,1,6))

print('goalBoard:\n')
printBoard(goalBoard)
print('btest1:\n')
printBoard(btest1)
print('btest2:\n')
printBoard(btest2)
print('(row,col) for 1 in btest1: ',(row(1,btest1),col(1,btest1)))
print('(row,col) for 6 in btest2: ',(row(6,btest2),col(6,btest2)))

print('\ntilesOutOfPlace(goalBoard,goalBoard):',tilesOutOfPlace(goalBoard,goalBoard) )

```

```

print('manhattanDistance(goalBoard,goalBoard):',manhattanDistance(goalBoard,goalBoard) )
print('tilesOutOfPlace(btest1,goalBoard):',tilesOutOfPlace(btest1,goalBoard) )
print('manhattanDistance(btest1,goalBoard):',manhattanDistance(btest1,goalBoard) )
print('tilesOutOfPlace(btest2,goalBoard):',tilesOutOfPlace(btest2,goalBoard) )
print('manhattanDistance(btest2,goalBoard):',manhattanDistance(btest2,goalBoard) )

print("\ngetChildren(goalBoard):")
for c in getChildren(goalBoard):
    printBoard(c)

print("\ngetChildren(btest1):")
for c in getChildren(btest1):
    printBoard(c)

print("\ngetChildren(btest2):")
for c in getChildren(btest2):
    printBoard(c)

```

goalBoard:

```

0 1 2
3 4 5
6 7 8

```

btest1:

```

3 1 2
4 5 0
6 7 8

```

btest2:

```

3 7 4
8 0 2
5 1 6

```

```

(row,col) for 1 in btest1: (0, 1)
(row,col) for 6 in btest2: (2, 2)

```

```

tilesOutOfPlace(goalBoard,goalBoard): 0
manhattanDistance(goalBoard,goalBoard): 0
tilesOutOfPlace(btest1,goalBoard): 4
manhattanDistance(btest1,goalBoard): 6
tilesOutOfPlace(btest2,goalBoard): 9
manhattanDistance(btest2,goalBoard): 18

```

getChildren(goalBoard):

```

3 1 2
0 4 5
6 7 8

```

```

1 0 2
3 4 5
6 7 8

```

getChildren(btest1):

```

3 1 0
4 5 2
6 7 8

```

```

3 1 2
4 5 8
6 7 0

```

```

3 1 2
4 0 5
6 7 8

```

getChildren(btest2):

```

3 0 4
8 7 2
5 1 6

```

```

3 7 4

```

```

8 1 2
5 0 6

3 7 4
0 8 2
5 1 6

3 7 4
8 2 0
5 1 6

```

Part B (30 pts)

Now you must write the main solver, following the template below.

Your main data structures and algorithm are as follows:

- **Open** : A priority queue implemented as a heap, ordered by $f(x)$; the elements of **Open** are tuples of the form $(f(B), B)$ for a board B . The heap will use the first element of the tuple to order the tuples (it is not used after it is popped from the queue).
 - At initialization, **Open** should contain only $(h(\text{start}), \text{start})$ (since $g(\text{start}) = 0$)
 - At each step, a minimal element (k, M) is popped off **OPEN** and processed ($k = f(M)$ is not further used):
 - Every child C of M (a board reachable in one move from M) is considered:
 - If C is not in **Closed** :
 - Add it to **Closed** with $g(C) = g(M) + 1$ and $\text{pred}(C) = M$ (since M is already in **Closed** you can get the distance $g(M)$ from there);
 - Add it to **Open** with $f(C) = g(C) + h(C)$
 - If C is already in **Closed**, determine whether the C just popped is closer to **start** (i.e., you found a shorter path to C):
 - If it is, update **Closed**[C] with the new values just described;
 - If not, ignore it
 - If C is the goal board, report success, print out the path if `trace = True`, and terminate.
 - If **Open** is empty, declare that the problem can not be solved and report the number of steps (should be close to 1814340).
 - **Closed** : Stores boards (instead of letters shown in the example above) and $g(\dots)$ values (distance from **start**).
 - At initialization, **Closed** should contain only **start** with $g(\text{start}) = 0$.
 - Every time you generate a child board from the board popped off **Open** you will also make sure the best possible path back to **start** is added or updated to **Closed**.

Be very clear that **Open** uses $f(\dots)$ as key and **Closed** stores only the distance $g(\dots)$; **Closed** stores all boards that have been generated at any time (so every board in **Open** is also in **Closed**).

Take careful note of the following points:

- You must use the function template with all parameters as listed:
 - `start` = the starting board
 - `goal` = the goal board (you will not have to use any other board for this problem, but should provide the ability to use other goals)
 - `heuristic` = the $h(\dots)$ heuristic function, one of the three given in the previous template
 - `Astar` = a flag to determine whether to use the A* algorithm:
 - if `Astar == true`, then the **OPEN** priority queue should be ordered by $f(x) = g(x) + h(x)$ as described in lecture on W 2/8;
 - if `Astar == false`, then **OPEN** should be ordered by $f(x) = h(x)$, i.e., ignore the distance $g(x)$ from the start state; this will give you heuristic (best-first) search, without ensuring that the solution path will be of minimum length. You will update the **Closed** list with the length of paths, as described above, but simply ignore this information.
 - `limit` = upper bound on the number of steps (pops from the **OPEN** list); since there are only $9!/2 = 1814340$ possible unique board reachable from any given start, the default limit should never be reached. You may of course

set it lower for testing.

- `trace` = flag to determine whether you should print out entire path from start to goal.

To summarize, here is what should happen under various parameter choices:

Heuristic	Astar	Behavior
noHeuristic	True	Breadth-First Search, but storing shortest paths
noHeuristic	False	Breadth-First Search but ignoring duplicates
tilesOutOfPlace	True	A* with tOOP heuristic
tilesOutOfPlace	False	Best-first search with tOOP
manhattanDistance	True	A* with nD heuristic
manhattanDistance	False	Best-first search with mD

```
In [24]: def EightPuzzle(start,goal=goalBoard,heuristic=manhattanDistance,Astar=True,limit=200000,trace=False):
pass      # your code here
```

You must test each of the six possible strategies given in the table above on the examples 0 - 7 given below; in addition, you must test the `manhattanDistance` and A* strategy on example 8 (which is not solvable).

You should get something very similar to what is shown here; because the order in which you generate children board may not be exactly the same as mine, there may be some differences, but the overall patterns should be similar.

```
In [2]: puzzle = [None]*10 # shortest solution
puzzle[0] = ((1, 0, 2), (3, 4, 5), (6, 7, 8)) # 1 step
puzzle[1] = ((1,4,2), (5,0,8), (3,6,7)) # 8 steps
puzzle[2] = ((1,5,2), (3,0,8), (4,6,7)) # 10 steps
puzzle[3] = ((0,5,1), (2,3,8), (6,4,7)) # 16 steps
puzzle[4] = ((3,1,4), (8,7,2), (5,0,6)) # 19 steps
puzzle[5] = ((1,0,3), (4,2,6), (7,5,8)) # 21 steps
puzzle[6] = ((2,1,6), (5,3,0), (4,7,8)) # 23/25 steps
puzzle[7] = ((8,7,6), (5,4,3), (2,1,0)) # 28 steps
puzzle[8] = ((0, 1, 5), (3, 4, 2), (6, 7, 8)) # no solution after ~181440 = 9!/2 steps
```

```
----- Puzzle(0) -----
No heuristic:      Solution of length 1 found ( 1 step ).
Tiles out of place:  Solution of length 1 found ( 1 step ).
Manhattan distance: Solution of length 1 found ( 1 step ).
No Heuristic:, A*: Solution of length 1 found ( 1 step ).
Tiles out of place, A*: Solution of length 1 found ( 1 step ).
Manhattan distance, A*: Solution of length 1 found ( 1 step ).
----- Puzzle(1) -----
No heuristic:      Solution of length 62 found ( 940 steps ).
Tiles out of place:  Solution of length 40 found ( 110 steps ).
Manhattan distance: Solution of length 8 found ( 18 steps ).
No Heuristic:, A*: Solution of length 8 found ( 130 steps ).
Tiles out of place, A*: Solution of length 8 found ( 11 steps ).
Manhattan distance, A*: Solution of length 8 found ( 16 steps ).
----- Puzzle(2) -----
No heuristic:      Solution of length 44 found ( 940 steps ).
Tiles out of place:  Solution of length 52 found ( 407 steps ).
Manhattan distance: Solution of length 16 found ( 22 steps ).
No Heuristic:, A*: Solution of length 10 found ( 441 steps ).
Tiles out of place, A*: Solution of length 10 found ( 48 steps ).
Manhattan distance, A*: Solution of length 10 found ( 52 steps ).
----- Puzzle(3) -----
No heuristic:      Solution of length 128 found ( 5225 steps ).
Tiles out of place:  Solution of length 40 found ( 79 steps ).
Manhattan distance: Solution of length 50 found ( 244 steps ).
No Heuristic:, A*: Solution of length 16 found ( 5730 steps ).
Tiles out of place, A*: Solution of length 16 found ( 494 steps ).
Manhattan distance, A*: Solution of length 16 found ( 160 steps ).
----- Puzzle(4) -----
```

```

No heuristic:           Solution of length 55 found ( 949 steps ).
Tiles out of place:    Solution of length 33 found ( 119 steps ).
Manhattan distance:    Solution of length 49 found ( 242 steps ).
No Heuristic:, A*:     Solution of length 19 found ( 19972 steps ).
Tiles out of place, A*: Solution of length 19 found ( 1488 steps ).
Manhattan distance, A*: Solution of length 19 found ( 243 steps ).
----- Puzzle(5) -----
No heuristic:           Solution of length 117 found ( 5214 steps ).
Tiles out of place:    Solution of length 79 found ( 471 steps ).
Manhattan distance:    Solution of length 39 found ( 80 steps ).
No Heuristic:, A*:     Solution of length 21 found ( 47215 steps ).
Tiles out of place, A*: Solution of length 21 found ( 5055 steps ).
Manhattan distance, A*: Solution of length 21 found ( 880 steps ).
----- Puzzle(6) -----
No heuristic:           Solution of length 135 found ( 5229 steps ).
Tiles out of place:    Solution of length 55 found ( 415 steps ).
Manhattan distance:    Solution of length 51 found ( 119 steps ).
No Heuristic:, A*:     Solution of length 23 found ( 84209 steps ).
Tiles out of place, A*: Solution of length 23 found ( 10898 steps ).
Manhattan distance, A*: Solution of length 23 found ( 2185 steps ).
----- Puzzle(7) -----
No heuristic:           Solution of length 84 found ( 955 steps ).
Tiles out of place:    Solution of length 28 found ( 28 steps ).
Manhattan distance:    Solution of length 60 found ( 211 steps ).
No Heuristic:, A*:     Solution of length 28 found ( 170274 steps ).
Tiles out of place, A*: Solution of length 28 found ( 51207 steps ).
Manhattan distance, A*: Solution of length 28 found ( 807 steps ).
----- Puzzle(8) -----
Manhattan distance, A*: Not found after 181440 steps!

```

Part C (5 pts)

Discuss what you found out about these strategies for the 8-puzzle problem. For example, all the A* strategies should have found a minimal solution of the same length for a given puzzle. However, all the other numbers show some patterns but also some surprising variations from what you might have expected. Describe what you see and if possible, what are the reasons for this behavior.